Querying Semantic Web Data with SPARQL (and SPARQL 1.1)

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Semantic Web

"The Semantic Web is an extension of the current web in which information is given well-defined meaning, better enabling computers and people to work in cooperation."
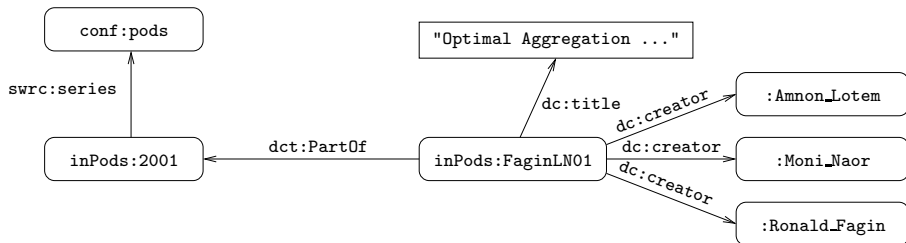
[Tim Berners-Lee et al. 2001.]

Specific goals:

- Build a description language with standard semantics
  - Make semantics machine-processable and understandable
- Incorporate logical infrastructure to reason about resources
- W3C proposals: Resource Description Framework (RDF) and SPARQL

An example of an RDF graph: DBLP

```
      : <http://dblp.l3s.de/d2r/resource/authors/>
   conf: <http://dblp.l3s.de/d2r/resource/conferences/>
  inPods: <http://dblp.l3s.de/d2r/resource/publications/conf/pods/>
   swrc: <http://swrc.ontoware.org/ontology#>
     dc: <http://purl.org/dc/elements/1.1/>
    dct: <http://purl.org/dc/terms/>
```

Querying RDF: SPARQL

- ▶ SPARQL is the W3C recommendation query language for RDF (January 2008).
    - ▶ SPARQL is a recursive acronym that stands for *SPARQL Protocol and RDF Query Language*

- ▶ SPARQL is a graph-matching query language.

- ▶ A SPARQL query consists of three parts:
    - ▶ Pattern matching: optional, union, filtering, . . .
    - ▶ Solution modifiers: projection, distinct, order, limit, offset, . . .
    - ▶ Output part: construction of new triples, ...

SPARQL in a nutshell

SPARQL in a nutshell

```
SELECT ?Author
```

SPARQL in a nutshell

```
SELECT ?Author
WHERE
{


}
```

SPARQL in a nutshell

```
SELECT ?Author
WHERE
{
  ?Paper      dc:creator      ?Author .
  ?Paper      dct:PartOf      ?Conf .
  ?Conf       swrc:series     conf:pods .
}
```

Outline of the talk

- ► RDF and SPARQL

- ► New features in SPARQL 1.1
  - ► Entailment regimes for RDFS and OWL
  - ► Navigational capabilities: Property paths
  - ► An operator to distribute the execution of a query

- ► Take-home message
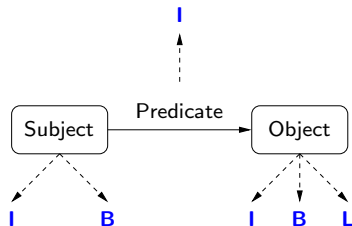
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RDF formal model
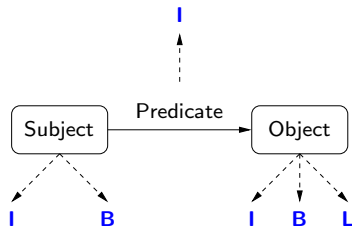
I	:	set of IRIs
B	:	set of blank nodes
L	:	set of literals

RDF formal model

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$(s, p, o) \in (\mathbf{I} \cup \mathbf{B}) \times \mathbf{I} \times (\mathbf{I} \cup \mathbf{B} \cup \mathbf{L})$ is called an RDF triple
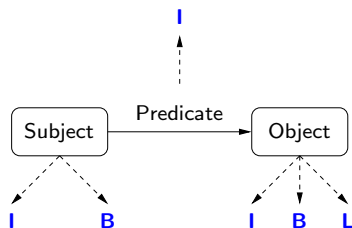
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L : set of literals

$(s, p, o) \in (\mathbf{I} \cup \mathbf{B}) \times \mathbf{I} \times (\mathbf{I} \cup \mathbf{B} \cup \mathbf{L})$ is called an RDF triple

A finite set of RDF triples is called an RDF graph

RDF formal model

Proviso

- We do not consider blank nodes in RDF graphs
  - $(s, p, o) \in \mathbf{I} \times \mathbf{I} \times (\mathbf{I} \cup \mathbf{L})$ is called an RDF triple

- We consider blank nodes in queries
  - Each blank node is assumed to start with _:, for example _:b and _:b_1

SPARQL: An algebraic syntax

V: set of variables

- Each variable is assumed to start with ?

Triple pattern: $t \in (\mathbf{I} \cup \mathbf{B} \cup \mathbf{V}) \times (\mathbf{I} \cup \mathbf{V}) \times (\mathbf{I} \cup \mathbf{B} \cup \mathbf{L} \cup \mathbf{V})$

- Examples: $(?X, \text{name}, \text{john})$, $(?X, \text{name}, ?Y)$, $(?X, \text{name}, \_\!:b)$

Basic graph pattern (bgp): Finite set of triple patterns

- Examples: $\{(?X, \text{knows}, ?Y), (?Y, \text{name}, \text{john})\}$,
  $\{(?X, \text{knows}, \_\!:b), (\_\!:b, \text{name}, \text{john})\}$

SPARQL: An algebraic syntax (cont'd)

Recursive definition of SPARQL graph patterns:

- ▶ Every basic graph pattern is a graph pattern

- ▶ If P_1, P_2 are graph patterns, then $(P_1$ AND $P_2)$, $(P_1$ OPT $P_2)$, $(P_1$ UNION $P_2)$ are graph pattern

- ▶ If P is a graph pattern and R is a *built-in condition*, then $(P$ FILTER $R)$ is a graph pattern

SPARQL query:

- ▶ If P is a graph pattern and W is a finite set of variables, then (SELECT W P) is a SPARQL query

Mappings: building block for the semantics

Definition

A mapping is a partial function:

$$\mu \;:\; \mathbf{V} \;\longrightarrow\; (\mathbf{I} \cup \mathbf{L})$$

The evaluation of a graph pattern results in a set of mappings.

Mappings: building block for the semantics

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A mapping is a partial function:

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Semantics of SPARQL: Basic graph patterns

Additional notation: $\sigma : \mathbf{B} \to (\mathbf{I} \cup \mathbf{L})$ is an instance mapping.

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Let P be a basic graph pattern

- $\mathrm{var}(P)$: set of variables mentioned in P

Definition

The evaluation of P over an RDF graph G, denoted by $[\![P]\!]_G$, is the set of mappings μ:

- $\mathrm{dom}(\mu) = \mathrm{var}(P)$
- there exists an instance mapping σ such that $\mu(\sigma(P)) \subseteq G$

Semantics of basic graph patterns: Some examples

Notation: t is used to represent $\{t\}$

Semantics of basic graph patterns: Some examples

Notation: t is used to represent $\{t\}$

graph	bgp	evaluation		
$(R_1, \text{name}, \text{john})$			?X	?Y
$(R_1, \text{email}, \text{J@ed.ex})$	$(?X, \text{name}, ?Y)$	μ_1:	R_1	john
$(R_2, \text{name}, \text{paul})$		μ_2:	R_2	paul

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$(R_2, \text{name}, \text{paul})$		

	?X	?Y
μ_1:	R_1	john
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graph	bgp	evaluation
$(R_1, \text{name}, \text{john})$		
$(R_1, \text{email}, \text{J@ed.ex})$	$(?X, \text{name}, \_{:}b)$	
$(R_2, \text{name}, \text{paul})$		

	?X
μ_3:	R_1
μ_4:	R_2

Compatible mappings

Definition

Mappings μ_1 and μ_2 are compatible if they agree in their common variables:

If $?X \in \mathrm{dom}(\mu_1) \cap \mathrm{dom}(\mu_2)$, then $\mu_1(?X) = \mu_2(?X)$.

Example

	?X	?Y	?Z	?V
μ_1 :	R_1	john		
μ_2 :	R_1		J@edu.ex	
μ_3 :			P@edu.ex	R_2

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$\mu_1 \cup \mu_2$:	R_1	john	J@edu.ex	

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$\mu_1 \cup \mu_2:$	R_1	john	J@edu.ex	
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	?X	?Y	?Z	?V
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$\mu_2:$	R_1		J@edu.ex	
$\mu_3:$			P@edu.ex	R_2
$\mu_1 \cup \mu_2:$	R_1	john	J@edu.ex	
$\mu_1 \cup \mu_3:$	R_1	john	P@edu.ex	R_2

- μ_2 and μ_3 are not compatible

Sets of mappings and operations

Let Ω_1 and Ω_2 be sets of mappings.

> **Definition**
>
> Join: extends mappings in Ω_1 with compatible mappings in Ω_2
>
> - $\Omega_1 \bowtie \Omega_2 = \{\mu_1 \cup \mu_2 \mid \mu_1 \in \Omega_1,\ \mu_2 \in \Omega_2 \text{ and } \mu_1,\ \mu_2 \text{ are compatible}\}$
>
> Difference: selects mappings in Ω_1 that cannot be extended with mappings in Ω_2
>
> - $\Omega_1 \smallsetminus \Omega_2 = \{\mu_1 \in \Omega_1 \mid \text{there is no mapping in } \Omega_2 \text{ compatible with } \mu_1\}$

Sets of mappings and operations

Definition

Union: includes mappings in Ω_1 and in Ω_2

- $\Omega_1 \cup \Omega_2 = \{\mu \mid \mu \in \Omega_1 \text{ or } \mu \in \Omega_2\}$

Left Outer Join: extends mappings in Ω_1 with compatible mappings in Ω_2 **if possible**

- $\Omega_1 \bowtie \Omega_2 = (\Omega_1 \bowtie \Omega_2) \cup (\Omega_1 \smallsetminus \Omega_2)$

Semantics of SPARQL: AND, UNION, OPT and SELECT

Given an RDF graph G

Definition

$\llbracket (P_1 \text{ AND } P_2) \rrbracket_G \quad =$

$\llbracket (P_1 \text{ UNION } P_2) \rrbracket_G \quad =$

$\llbracket (P_1 \text{ OPT } P_2) \rrbracket_G \quad =$

$\llbracket (\text{SELECT } W \ P) \rrbracket_G \quad =$

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$$\llbracket (P_1 \text{ OPT } P_2) \rrbracket_G \quad = \quad \llbracket P_1 \rrbracket_G \mathbin{⟗} \llbracket P_2 \rrbracket_G$$

$$\llbracket (\text{SELECT } W \ P) \rrbracket_G \quad = \quad \{ \mu_{|_W} \mid \mu \in \llbracket P \rrbracket_G \}$$

Semantics of SPARQL: AND, UNION, OPT and SELECT

Given an RDF graph G

Definition

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$$[\![(\text{SELECT } W \ P)]\!]_G \quad = \quad \{\mu_{|_W} \mid \mu \in [\![P]\!]_G\}$$

$$\text{dom}(\mu_{|_W}) = \text{dom}(\mu) \cap W \text{ and}$$
$$\mu_{|_W}(?X) = \mu(?X) \text{ for every } ?X \in \text{dom}(\mu_{|_W})$$

Semantics of SPARQL: An example

Example

$(R_1, \text{name}, \text{john})$
$(R_1, \text{email}, \text{J@ed.ex})$
$(R_2, \text{name}, \text{paul})$

$((?X, \text{name}, ?Y) \text{ OPT } (?X, \text{email}, ?E))$

Semantics of SPARQL: An example

> **Example**
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>
> $$((?X, \text{name}, ?Y) \text{ OPT } (?X, \text{email}, ?E))$$

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$$(\;(?X, \text{ name, } ?Y) \text{ OPT } (?X, \text{ email, } ?E)\;)$$

$?X$	$?Y$
R_1	john
R_2	paul

Semantics of SPARQL: An example

$(R_1, \text{name}, \text{john})$
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$(\,(?X, \text{name}, ?Y)\ \text{OPT}\ (?X, \text{email}, ?E)\,)$

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$?X$	$?Y$
R_1	john
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$?X$	$?E$
R_1	J@ed.ex

Semantics of SPARQL: An example

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Semantics of SPARQL: An example

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R_2	paul

$?X$	$?Y$	$?E$
R_1	john	J@ed.ex

$?X$	$?E$
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▶ from the Join

Semantics of SPARQL: An example

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$$(R_1, \text{email}, \text{J@ed.ex})$$
$$(R_2, \text{name}, \text{paul})$$

$$((?X, \text{name}, ?Y) \ \text{OPT} \ (?X, \text{email}, ?E))$$

$?X$	$?Y$
R_1	john
R_2	paul

$?X$	$?Y$	$?E$
R_2	paul	

$?X$	$?E$
R_1	J@ed.ex

- from the Difference

Semantics of SPARQL: An example

Example

$$(R_1, \text{name}, \text{john})$$
$$(R_1, \text{email}, \text{J@ed.ex})$$
$$(R_2, \text{name}, \text{paul})$$

$$((?X, \text{name}, ?Y) \text{ OPT } (?X, \text{email}, ?E))$$

$?X$	$?Y$
R_1	john
R_2	paul

$?X$	$?Y$	$?E$
R_1	john	J@ed.ex
R_2	paul	

$?X$	$?E$
R_1	J@ed.ex

- from the Union

Filter expressions (value constraints)

Filter expression: $(P \text{ FILTER } R)$

- P is a graph pattern
- R is a built-in condition

We consider in R:

- equality $=$ among variables and RDF terms
- unary predicate bound
- boolean combinations $(\wedge, \vee, \neg)$

A mapping μ satisfies a condition R ($\mu \models R$) if:

Satisfaction of value constraints

A mapping μ satisfies a condition R ($\mu \models R$) if:

- R is $?X = c$, $?X \in \text{dom}(\mu)$ and $\mu(?X) = c$
- R is $?X = ?Y$, $?X, ?Y \in \text{dom}(\mu)$ and $\mu(?X) = \mu(?Y)$
- R is $\text{bound}(?X)$ and $?X \in \text{dom}(\mu)$

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- usual rules for Boolean connectives

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- R is $\text{bound}(?X)$ and $?X \in \text{dom}(\mu)$
- usual rules for Boolean connectives

Definition

FILTER : selects mappings that satisfy a condition

$$[\![(P \text{ FILTER } R)]\!]_G \;\; = \;\; \{\mu \in [\![P]\!]_G \mid \mu \models R\}$$

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- New features in SPARQL 1.1
    - Entailment regimes for RDFS and OWL
    - Navigational capabilities: Property paths
    - An operator to distribute the execution of a query

- Take-home message

SPARQL 1.1

A new version of SPARQL was released in March 2013: SPARQL 1.1

Some new features in SPARQL 1.1:

- ▶ Entailment regimes for RDFS and OWL

- ▶ Navigational capabilities: Property paths

- ▶ An operator (SERVICE) to distribute the execution of a query

Also in this version: Nesting of SELECT expressions, aggregates and some forms of negation (NOT EXISTS, MINUS)

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  - An operator to distribute the execution of a query

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Syntax of RDFS

RDFS extends RDF with a schema vocabulary: subPropertyOf (rdfs:sp), subClassOf (rdfs:sc), domain (rdfs:dom), range (rdfs:range), type (rdfs:type).

Syntax of RDFS

RDFS extends RDF with a schema vocabulary: subPropertyOf (rdfs:sp), subClassOf (rdfs:sc), domain (rdfs:dom), range (rdfs:range), type (rdfs:type).

How do we evaluate a query over RDFS data?

A simple SPARQL query: (Messi, rdfs:type, person)

Semantics of RDFS
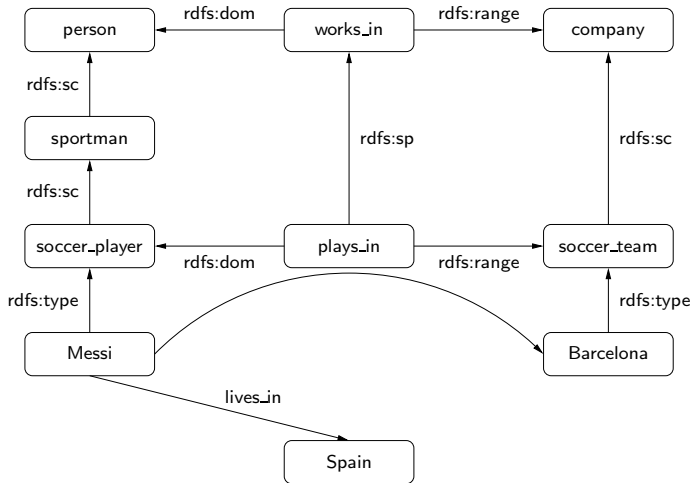
Checking whether a triple t is in a graph G is the basic step when answering queries over RDF.

- For the case of RDFS, we need to check whether t is implied by G

The notion of entailment in RDFS can be defined as for first-order logic.

This notion can also be characterized by a set of inference rules.

An inference system for RDFS

Sub-property :
$$\frac{(\mathcal{A},\ \text{rdfs:sp},\ \mathcal{B})\ (\mathcal{B},\ \text{rdfs:sp},\ \mathcal{C})}{(\mathcal{A},\ \text{rdfs:sp},\ \mathcal{C})}$$

$$\frac{(\mathcal{A},\ \text{rdfs:sp},\ \mathcal{B})\ (\mathcal{X},\ \mathcal{A},\ \mathcal{Y})}{(\mathcal{X},\ \mathcal{B},\ \mathcal{Y})}$$

Subclass :
$$\frac{(\mathcal{A},\ \text{rdfs:sc},\ \mathcal{B})\ (\mathcal{B},\ \text{rdfs:sc},\ \mathcal{C})}{(\mathcal{A},\ \text{rdfs:sc},\ \mathcal{C})}$$

$$\frac{(\mathcal{A},\ \text{rdfs:sc},\ \mathcal{B})\ (\mathcal{X},\ \text{rdfs:type},\ \mathcal{A})}{(\mathcal{X},\ \text{rdfs:type},\ \mathcal{B})}$$

Typing :
$$\frac{(\mathcal{A},\ \text{rdfs:dom},\ \mathcal{B})\ (\mathcal{X},\ \mathcal{A},\ \mathcal{Y})}{(\mathcal{X},\ \text{rdfs:type},\ \mathcal{B})}$$

$$\frac{(\mathcal{A},\ \text{rdfs:range},\ \mathcal{B})\ (\mathcal{X},\ \mathcal{A},\ \mathcal{Y})}{(\mathcal{Y},\ \text{rdfs:type},\ \mathcal{B})}$$

Theorem (H03,MPG09,GHM11)

The previous system of inference rules characterize the notion of entailment in RDFS (without blank nodes).

Thus, a triple t can be deduced from an RDF graph G ($G \models t$) iff t can be deduced from G by applying the inference rules a finite number of times.

An entailment regime for RDFS in SPARQL 1.1

Basic graph patterns are evaluated by considering RDFS entailment.

Definition

The evaluation of a bgp P over an RDF graph G, denoted by $[\![P]\!]_G$, is the set of mappings μ:

- $\mathrm{dom}(\mu) = \mathrm{var}(P)$

- there exists an instance mapping σ such that for every $t \in P$:
  $G \models \mu(\sigma(t))$

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The semantics of AND, UNION, OPT, FILTER and SELECT are defined as before.

- RDFS entailment is only used at the level of bgps

Entailment regimes in SPARQL 1.1: Some observations

- ▶ SPARQL 1.1 can be used to query not only data but also schema information
  - ▶ For example: $(?X, \text{rdfs:sc}, \text{person})$

Entailment regimes in SPARQL 1.1: Some observations

- SPARQL 1.1 can be used to query not only data but also schema information
  - For example: $(?X, \text{rdfs:sc}, \text{person})$

- Basic graph patterns can also be evaluated by considering OWL entailment.
  - $G \models \mu(\sigma(t))$ has to be defined according to the semantics of OWL

▶ What are the consequences of considering entailment only at the level bgps?

Example

Let G be a graph consisting of (john, rdf:type, student) together with:

$$\left.\begin{array}{l} (\text{student}, \text{rdfs:sc}, u) \\ (u, \text{owl:union}, l) \\ (l, \text{rdf:first}, \text{undergrad}) \\ (l, \text{rdf:rest}, r) \\ (r, \text{rdf:first}, \text{grad}) \\ (r, \text{rdf:rest}, \text{rdf:nil}) \end{array}\right\} \text{axiom student} \sqsubseteq (\text{undergrad} \sqcup \text{grad})$$

What should be the answer to
$P = ((?X, \text{rdfs:type}, \text{undergrad}) \text{ UNION } (?X, \text{rdfs:type}, \text{grad}))$?

▶ Under the current semantics: $[\![P]\!]_G = \emptyset$

- It is possible to define a certain-answers semantics for SPARQL 1.1.
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  But what happens if we focus on the case of RDFS?
  - ▶ The semantics do not coincide as the following operator can be expressed in the language:

$$[\![(P_1 \text{ MINUS } P_2)]\!]_G \;\; = \;\; [\![P_1]\!]_G \smallsetminus [\![P_2]\!]_G$$

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Open issues

- ▶ How natural is the semantics of SPARQL 1.1? Is it a good semantics? Why?
- ▶ Under which (natural) restrictions these two semantics coincide?

Outline of the talk

- RDF and SPARQL

- New features in SPARQL 1.1
  - Entailment regimes for RDFS and OWL
  - Navigational capabilities: Property paths
  - An operator to distribute the execution of a query

- Take-home message

SPARQL provides limited navigational capabilities
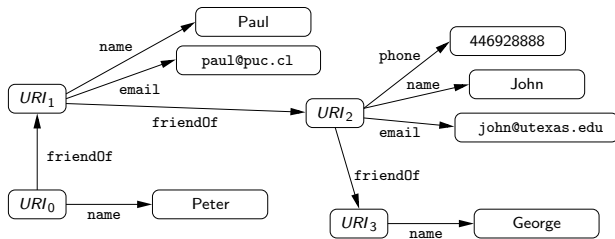
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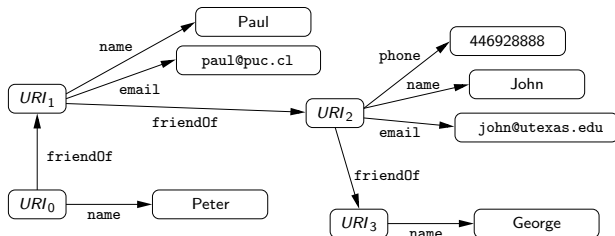
(SELECT ?X ((?X, friendOf, ?Y) AND (?Y, name, George)))

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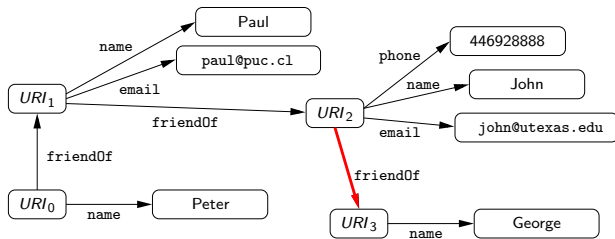
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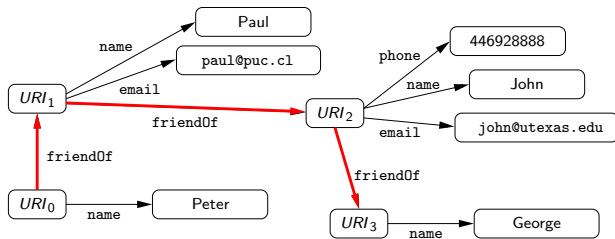
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A possible solution: Property paths
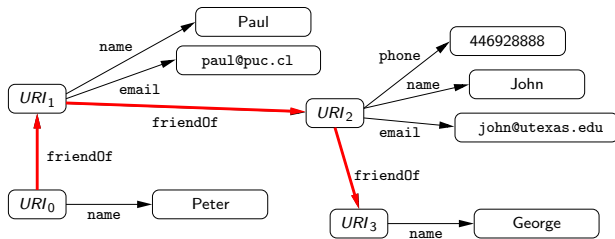
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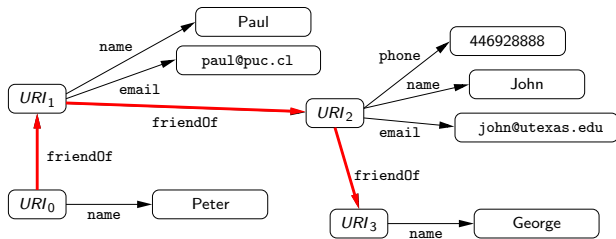
(SELECT ?X ((?X, (friendOf)*, ?Y) AND (?Y, name, George)))

Syntax of property paths:

$$exp \ := \ a \ | \ exp/exp \ | \ exp|exp \ | \ exp^*$$

where $a \in \mathbf{I}$

Syntax of property paths:

$$exp \quad := \quad a \mid exp/exp \mid exp|exp \mid exp^*$$

where $a \in \mathbf{I}$

Other expressions are allowed:

$\hat{}\,exp$	:	inverse path
$!(a_1|\ldots|a_n)$	:	an IRI which is not one of a_i ($1 \leq i \leq n$)

Evaluating property paths

The evaluation of a property path over an RDF graph G is defined as follows:

Evaluating property paths

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$$
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[\![a]\!]_G &= \{(x, y) \mid (x, a, y) \in G\} \\
[\![exp_1/exp_2]\!]_G &= \{(x, y) \mid \exists z\ (x, z) \in [\![exp_1]\!]_G \text{ and} \\
&\qquad\qquad\qquad\quad (z, y) \in [\![exp_2]\!]_G\}
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[\![exp_1|exp_2]\!]_G &= [\![exp_1]\!]_G \cup [\![exp_2]\!]_G \\
[\![exp^*]\!]_G &= \{(a,a) \mid a \text{ is an IRI in } G\} \cup [\![exp]\!]_G \cup \\
&\qquad [\![exp/exp]\!]_G \cup [\![exp/exp/exp]\!]_G \cup \cdots
\end{aligned}
$$

Property paths in SPARQL 1.1

New element in SPARQL 1.1: A triple of the form (x, exp, y)

- exp is a property path
- x (resp. y) is either an element from **I** or a variable

Property paths in SPARQL 1.1

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- x (resp. y) is either an element from $\mathbf{I}$ or a variable

Example

- $(?X, (\text{rdfs:sc})^*, \text{person})$: Verifies whether the value stored in $?X$ is a subclass of person

- $(?X, (\text{rdfs:sp})^*, ?Y)$: Verifies whether the value stored in $?X$ is a subproperty of the value stored in $?Y$

Semantics of property paths

Evaluation of $t = (?X, exp, ?Y)$ over an RDF graph G is the set of mappings μ such that:

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Other cases are defined analogously.

Example

- $((?X, \text{KLM}/(\text{KLM})^*, ?Y)$ FILTER $\neg(?X =?Y))$: It is possible to go from $?X$ to $?Y$ by using the airline KLM, where $?X$, $?Y$ are different cities

List the pairs *a*, *b* of cities such that there is a way to travel from *a* to *b*.

SPARQL 1.1: Entailment regimes and property paths

List the pairs a, b of cities such that there is a way to travel from a to b.
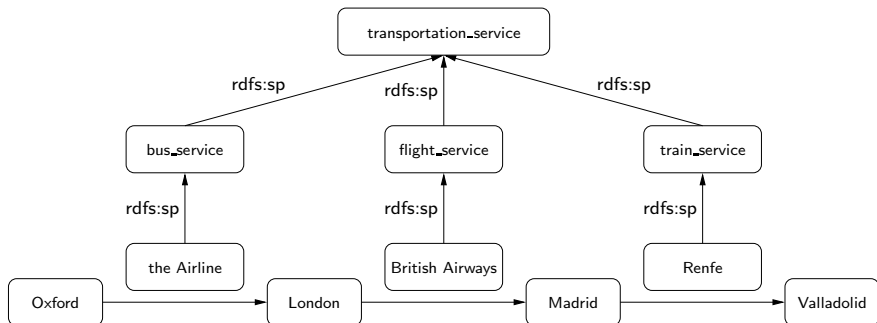
In SPARQL 1.1: $(?X, \text{transportation\_service}^*, ?Y)$

▶ Previous query can be expressed in SPARQL 1.1 as the intermediate form of navigation involves RDFS vocabulary.

Not expressible: List pairs a, b of persons that are connected through a path of nodes certified by certifying_agency [RK13]:
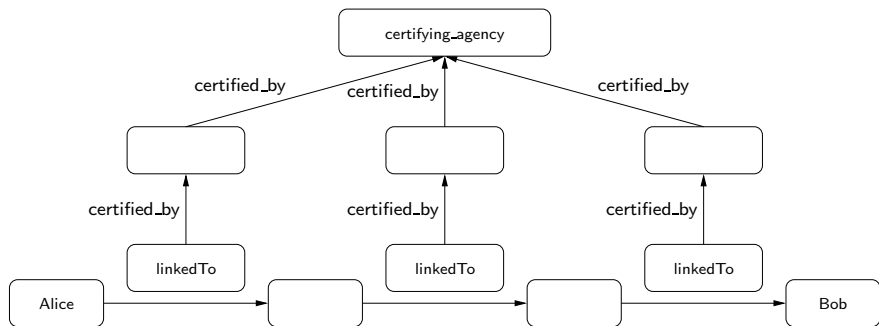
Navigational capabilities in SPARQL 1.1: Some observations (cont'd)

- ► Some proposals solve the aforementioned issues: nSPARQL [PAG10], nested monadically defined queries [RK13], triple algebra [LRV13]
  - ► RDFS entailment can be handled in these proposals by using navigational capabilities

- ▶ Some proposals solve the aforementioned issues: nSPARQL [PAG10], nested monadically defined queries [RK13], triple algebra [LRV13]
    - ▶ RDFS entailment can be handled in these proposals by using navigational capabilities

Open issues

- ▶ How can OWL entailment be handled in these proposals?
- ▶ What navigational capabilities should be added to SPARQL 1.1?
- ▶ There is a need for query languages that can return paths

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- Take-home message

RFD graphs can be interconnected

```
        : <http://dblp.l3s.de/d2r/resource/authors/>
 dbpedia: <http://dbpedia.org/resource/>
     rdf: <http://www.w3.org/1999/02/22-rdf-syntax-ns#>
    rdfs: <http://www.w3.org/2000/01/rdf-schema#>
     owl: <http://www.w3.org/2002/07/owl#>
    yago: <http://dbpedia.org/class/yago>
     dbo: <http://dbpedia.org/ontology/>
```
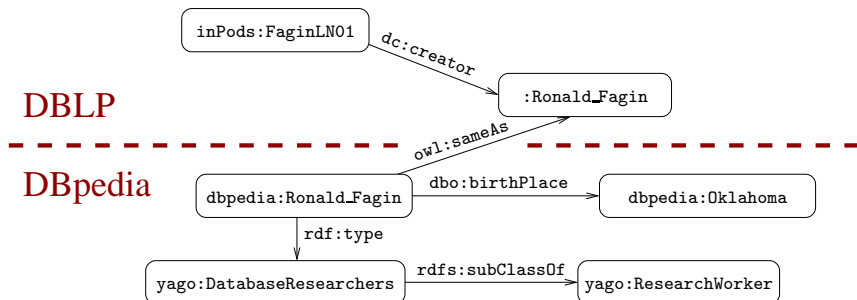
Querying interconnected RDF graphs

Retrieve the authors that have published in PODS and were born in Oklahoma:

```
SELECT ?Author
WHERE
{
  ?Paper        dc:creator      ?Author .
  ?Paper        dct:PartOf      ?Conf .
  ?Conf         swrc:series     conf:pods .
  SERVICE <http://dbpedia.org/sparql> {
    ?Person       owl:sameAs      ?Author .
    ?Person       dbo:birthPlace dbpedia:Oklahoma . }
}
```

Federation in SPARQL 1.1

New rule to generate graph patterns:

- If P is a graph pattern and $c \in (\mathbf{I} \cup \mathbf{V})$, then (SERVICE c P) is a graph pattern.

Federation in SPARQL 1.1

New rule to generate graph patterns:

- If P is a graph pattern and $c \in (\mathbf{I} \cup \mathbf{V})$, then (SERVICE c P) is a graph pattern.

We will define the semantics of this new operator.

- This corresponds with the official semantics for (SERVICE c P) with $c \in \mathbf{I}$

- (SERVICE $?X$ P) is allowed in the official specification of SPARQL 1.1, but its semantics is not defined

Semantics of SERVICE

ep(·): Partial function from **I** to the set of all RDF graphs

- ▶ If $c \in$ dom(ep), then ep(c) is the RDF graph associated with the endpoint accessible via c

Semantics of SERVICE

$ep(\cdot)$: Partial function from **I** to the set of all RDF graphs

- If $c \in \text{dom}(ep)$, then $ep(c)$ is the RDF graph associated with the endpoint accessible via c

Definition (BACP13)

The evaluation of $P = (\text{SERVICE } c \ P_1)$ over an RDF graph G is defined as:

- if $c \in \text{dom}(ep)$, then $[\![P]\!]_G = [\![P_1]\!]_{ep(c)}$
- if $c \in \textbf{I} \setminus \text{dom}(ep)$, then $[\![P]\!]_G = \{\mu_\emptyset\}$ (where $\mu_\emptyset$ is the mapping with empty domain)
- if $c \in \textbf{V}$, then

$$[\![P]\!]_G \quad = \quad \bigcup_{a \in \text{dom}(ep)} \left([\![P_1]\!]_{ep(a)} \bowtie \{\mu_{c \to a}\} \right),$$

where $\mu_{c \to a}$ is a mapping such that $\text{dom}(\mu_{c \to a}) = \{c\}$ and $\mu_{c \to a}(c) = a$

Are variables useful in SERVICE queries?

Consider the query:

$(?X, \text{service\_address}, ?Y)$ AND (SERVICE $?Y$ $(?N, \text{email}, ?E)$)

Are variables useful in SERVICE queries?

Consider the query:

$(?X, \text{service\_address}, ?Y)$ AND $(\text{SERVICE } ?Y \ (?N, \text{email}, ?E))$

There is a simple strategy to compute the answer to this query.

- ► Can this strategy be generalized?

How can we evaluate SERVICE queries?

We need some notion of boundedness

- A variable $?X$ is bound in a graph pattern P if for every RDF graph G and every $\mu \in \llbracket P \rrbracket_G$, it holds that $?X \in \text{dom}(\mu)$ and $\mu(?X)$ is mentioned in G

First attempt: Graph pattern P can be evaluated if for every sub-pattern (SERVICE $?X$ P_1) of P, it holds that $?X$ is bound in P

- $?Y$ is bound in
  $(?X, \text{service\_address}, ?Y)$ AND (SERVICE $?Y$ $(?N, \text{email}, ?E)$)

The first attempt: Too restrictive

Consider the query:

$(?X, \text{service\_description}, ?Z)$ UNION

$$\Big((?X, \text{service\_address}, ?Y) \text{ AND } (\text{SERVICE } ?Y \ (?N, \text{email}, ?E))\Big)$$

$?Y$ is not bound in this query, but there is a simple strategy to evaluate it.

The first attempt: Not appropriate for nested SERVICE operators

Consider the query:

$$(?U_1, \text{related\_with}, ?U_2) \ \text{AND}$$
$$\left[\text{SERVICE} \ ?U_1 \ \left((?N, \text{email}, ?E) \ \text{OPT} \right.\right.$$
$$\left.\left. (\text{SERVICE} \ ?U_2 \ (?N, \text{phone}, ?F)) \right) \right]$$

Solving the problems . . .

Notation: $\mathcal{T}(P)$ is the *parse* tree of P, in which every node corresponds to a sub-pattern of P

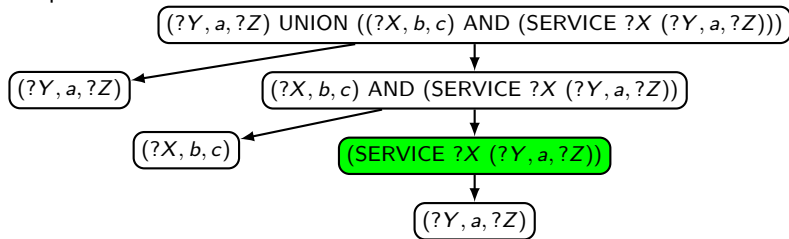Parse tree of $(?Y, a, ?Z)$ UNION $((?X, b, c)$ AND $($SERVICE $?X$ $(?Y, a, ?Z)))$:
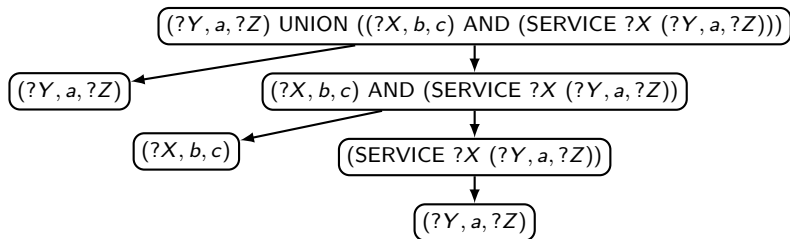
A more appropriate notion of boundedness

Definition (BACP13)

A graph pattern P is service-bound if for every node u of $\mathcal{T}(P)$ with label (SERVICE $?X$ P_1), it holds that:

- there exists a node v of $\mathcal{T}(P)$ with label P_2 such that v is an ancestor of u in $\mathcal{T}(P)$ and $?X$ is bound in P_2
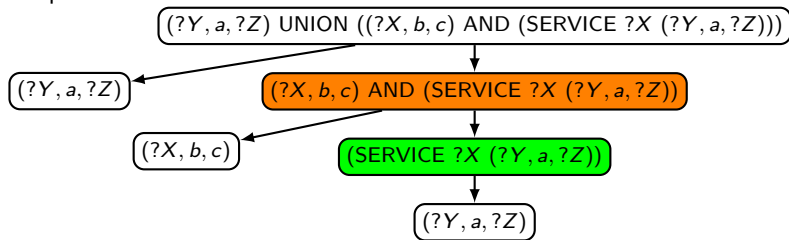- P_1 is service-bound

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Examples:

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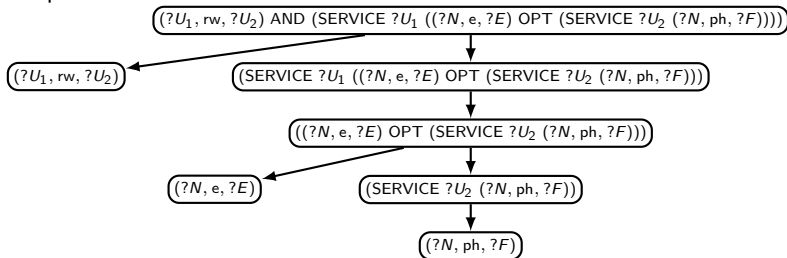- P_1 is service-bound

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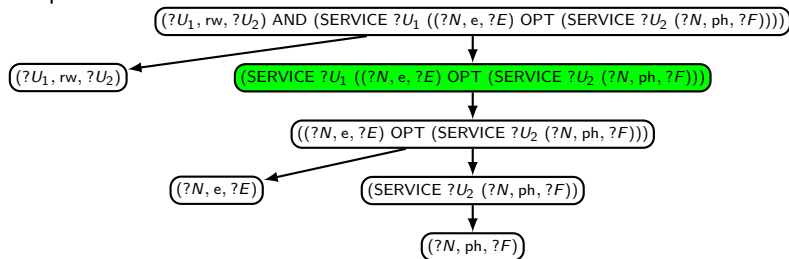
Examples:

$$\boxed{(?Y, a, ?Z)}$$

A more appropriate notion of boundedness

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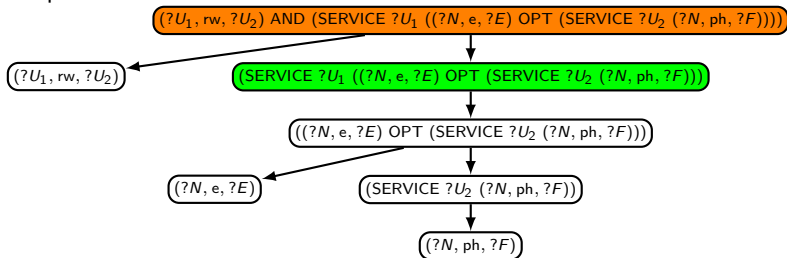- P_1 is service-bound

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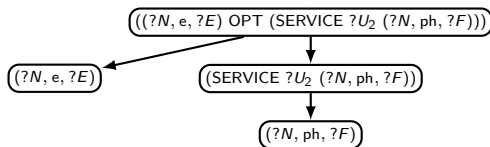- P_1 is service-bound

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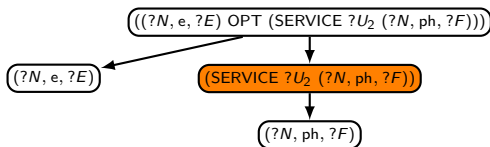- P_1 is service-bound

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- P_1 is service-bound

Examples:

A more appropriate notion of boundedness (cont'd)

But we still have a problem:

Proposition (BACP13)

The problem of verifying, given a graph pattern P, whether P is service-bound is undecidable.

We consider a (syntactic) sufficient condition for service-boundedness.

An appropriate notion: Service-safeness

The set of strongly bound variables in P, denoted by $SB(P)$, is recursively defined as follows:

- if P is a bgp, then $SB(P) = var(P)$
- if $P = (P_1 \text{ AND } P_2)$, then $SB(P) = SB(P_1) \cup SB(P_2)$
- if $P = (P_1 \text{ UNION } P_2)$, then $SB(P) = SB(P_1) \cap SB(P_2)$
- if $P = (P_1 \text{ OPT } P_2)$, then $SB(P) = SB(P_1)$
- if $P = (P_1 \text{ FILTER } R)$, then $SB(P) = SB(P_1)$
- if $P = (\text{SERVICE } c \ P_1)$, then $SB(P) = \emptyset$

Definition (BACP13)

A graph pattern P is service-safe if for every node u of $\mathcal{T}(P)$ with label (SERVICE $?X$ P_1), it holds that:

- there exists a node v of $\mathcal{T}(P)$ with label P_2 such that v is an ancestor of u in $\mathcal{T}(P)$ and $?X \in \text{SB}(P_2)$
- P_1 is service-safe

If P is service-safe, then there is a strategy to evaluate P without considering all possible SPARQL endpoints.

An appropriate notion: Service-safeness (cont'd)

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A graph pattern P is service-safe if for every node u of $\mathcal{T}(P)$ with label (SERVICE $?X$ P_1), it holds that:

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- P_1 is service-safe

If P is service-safe, then there is a strategy to evaluate P without considering all possible SPARQL endpoints.

Open issue

Is service-safeness the right condition to ensure that a query containing the SERVICE operator can be executed? Why?

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- Take-home message

Take-home message

- ▶ RDF is the framework proposed by the W3C to represent information in the Web

- ▶ SPARQL is the W3C recommendation query language for RDF (January 2008)

- ▶ SPARLQ 1.1 is the new version of SPARQL (March 2013)

- ▶ SPARQL 1.1 includes some interesting and useful new features
  - ▶ Entailment regimes for RDFS and OWL, navigational capabilities and an operator to distribute the execution of a query
  - ▶ There are some interesting open issues about these features

Thank you!

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